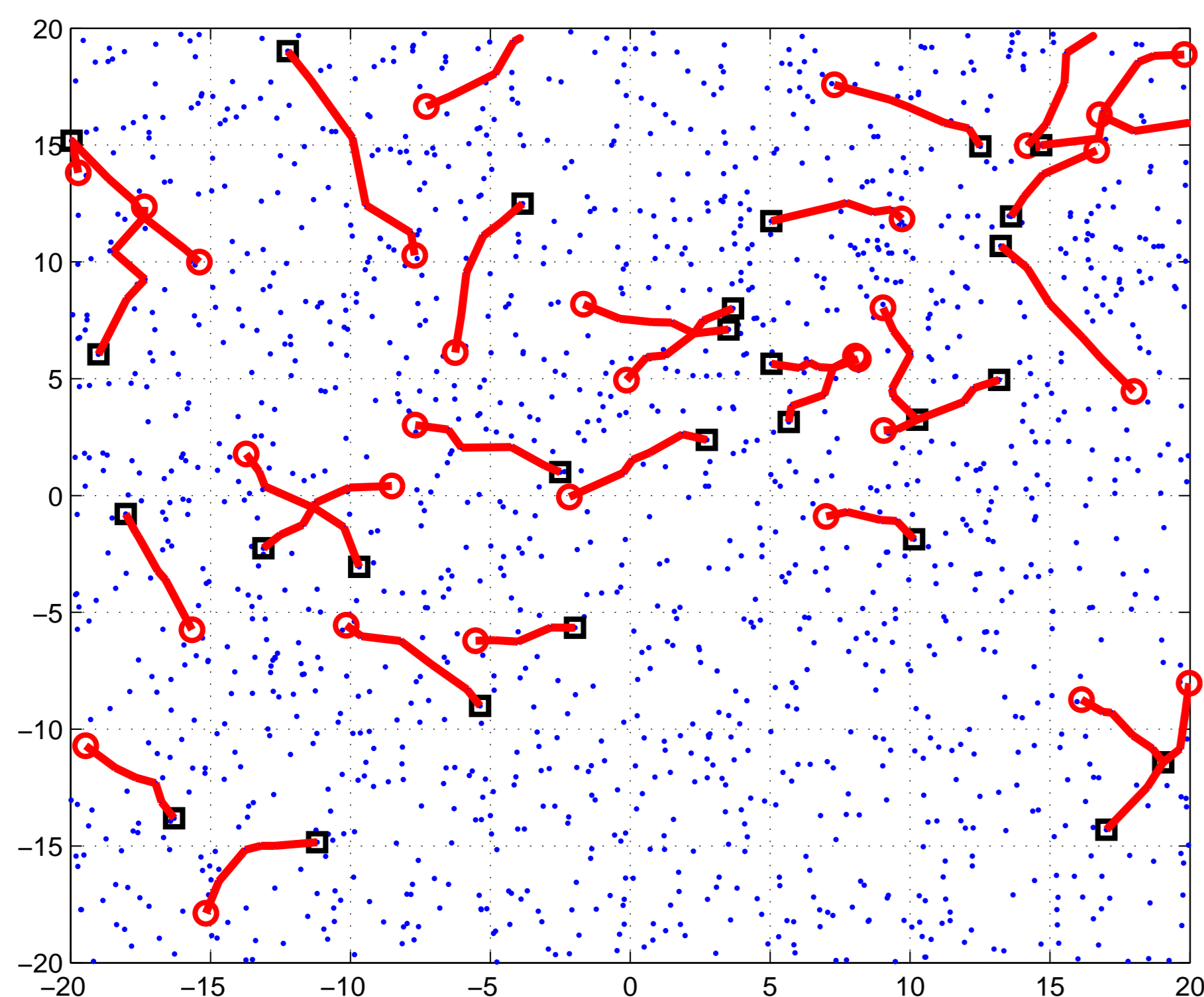


OVERVIEW

- Employ a combination of tools from two paradigms, **stochastic geometry** and **TASEP** (Totally Asymmetric Simple Exclusion Process) to analyze multihop wireless ad hoc networks.
- Provide valuable insights from a system design stand-point.

Prior work: single-hop analysis, backlogged nodes, simplified models.

SYSTEM MODEL



Sources \sim PPP (λ_s) ($= \delta\lambda$)
 Potential relays \sim PPP (λ_r) ($= (1-\delta)\lambda$).
 Channel: Rayleigh i.i.d., path loss exponent γ .

Routing Strategy:

In each flow, the next-hop node is the n^{th} -nearest-neighbor that lies within $\pm\phi/2$ of the axis to the destination.

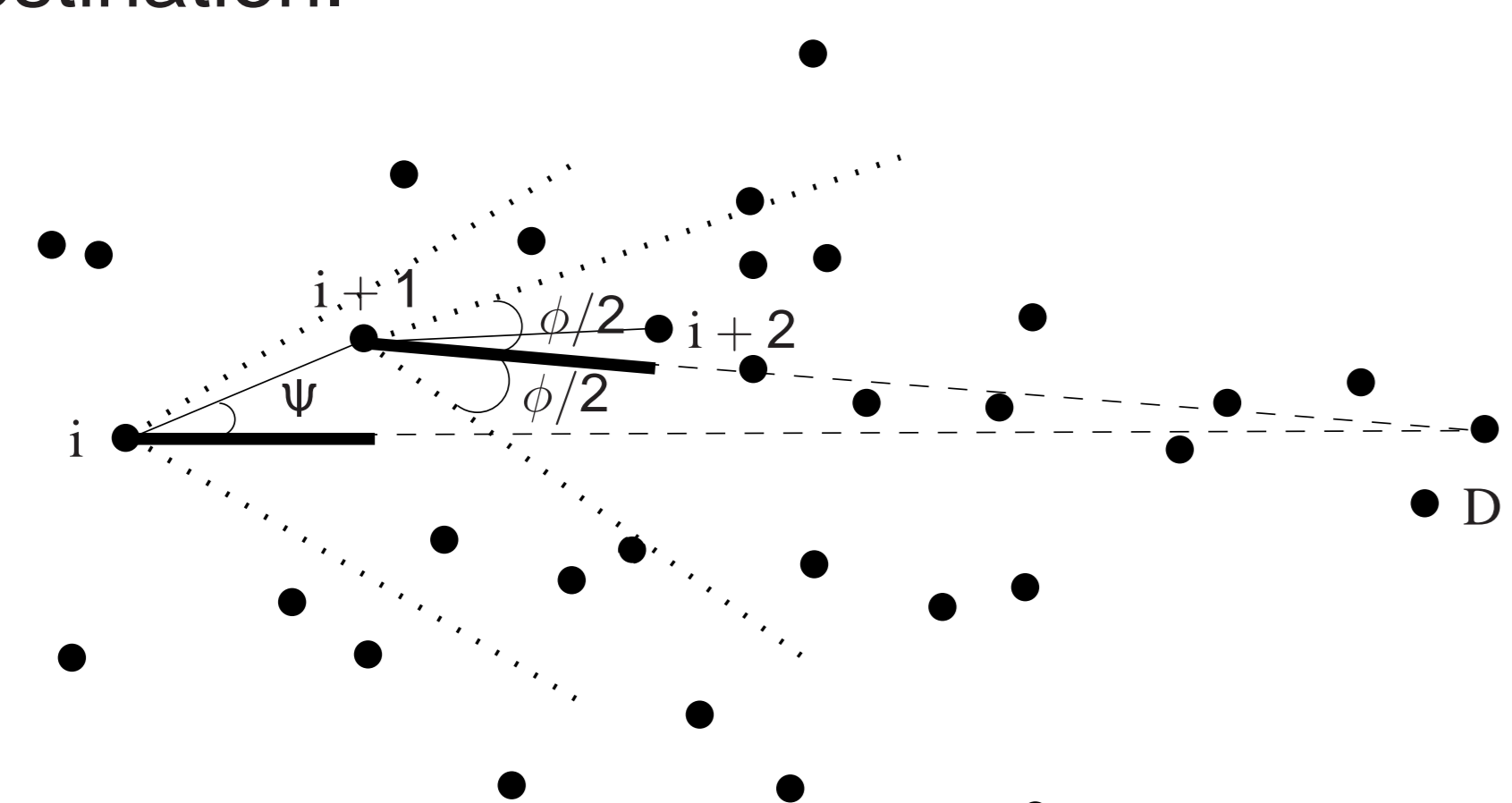


Illustration of nearest-neighbor ($n = 1$) routing.

MAC Schemes:

- Randomized TDMA (r-TDMA):** The transmitting node in each time slot is chosen uniformly randomly from the set of all nodes in the flow *having a packet*.
- ALOHA** with contention probability q .

Performance Metric:

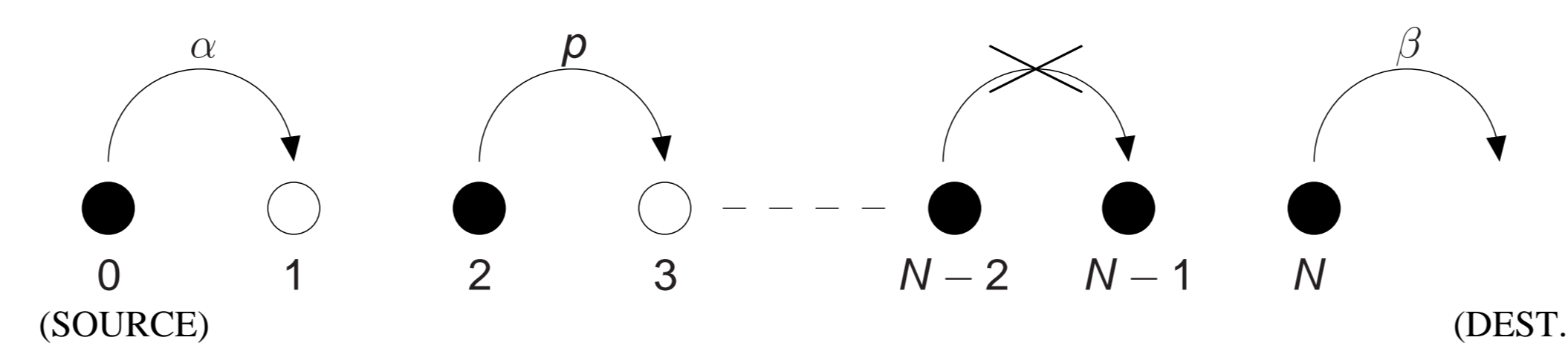
Transport Density is the average number of bit-meters successfully delivered per second per unit surface area: $\rho_{\text{trans}} = \delta TD$, where T is the throughput of a 'typical' flow. D denotes the progress of packets from source to destination.

THE TOTALLY ASYMMETRIC SIMPLE EXCLUSION PROCESS

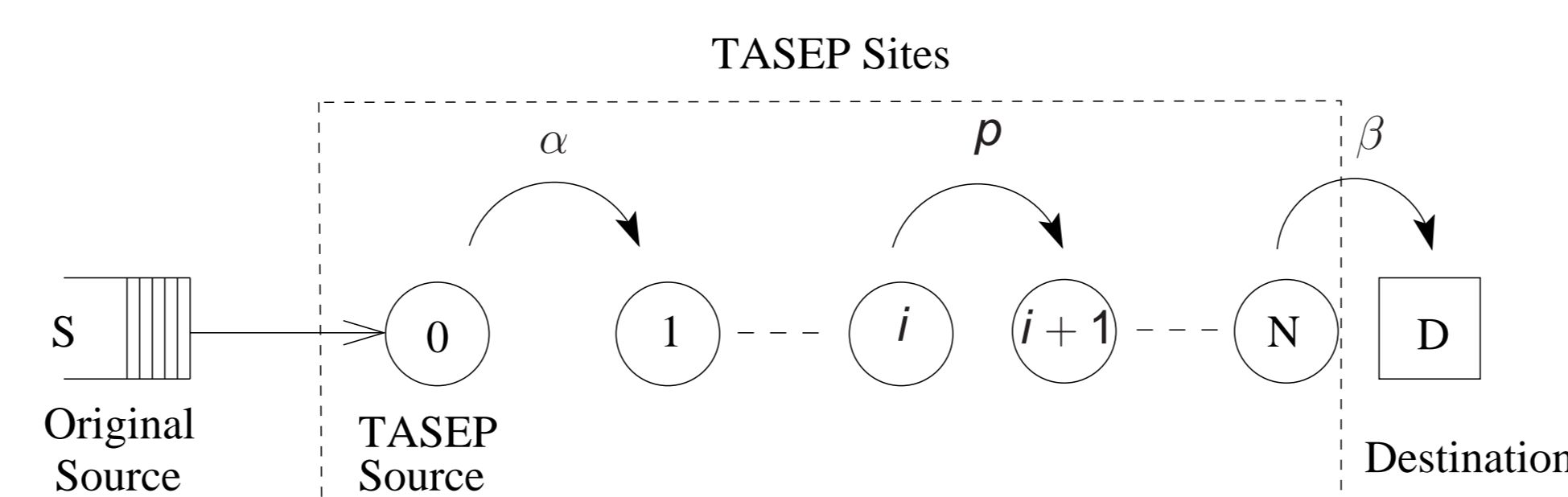
Transmission Policy:

- All the buffering in the network is performed at source nodes.
- Transmissions are not attempted if the adjacent relay's buffer already contains a packet.
- Packets are retransmitted until successfully received.

TASEP Model:



Network flows:



ANALYSIS OF R-TDMA-BASED WIRELESS NETWORKS

Proposition 1 For the r-TDMA-based ad hoc network, the probability of a successful transmission $p_s = \mathbb{P}[SIR > \beta]$ from any node to its n^{th} -nearest-neighbor in a sector ϕ is

$$p_s = \left(\frac{(1-\delta)\phi}{(1-\delta)\phi + 2\delta c} \right)^n,$$

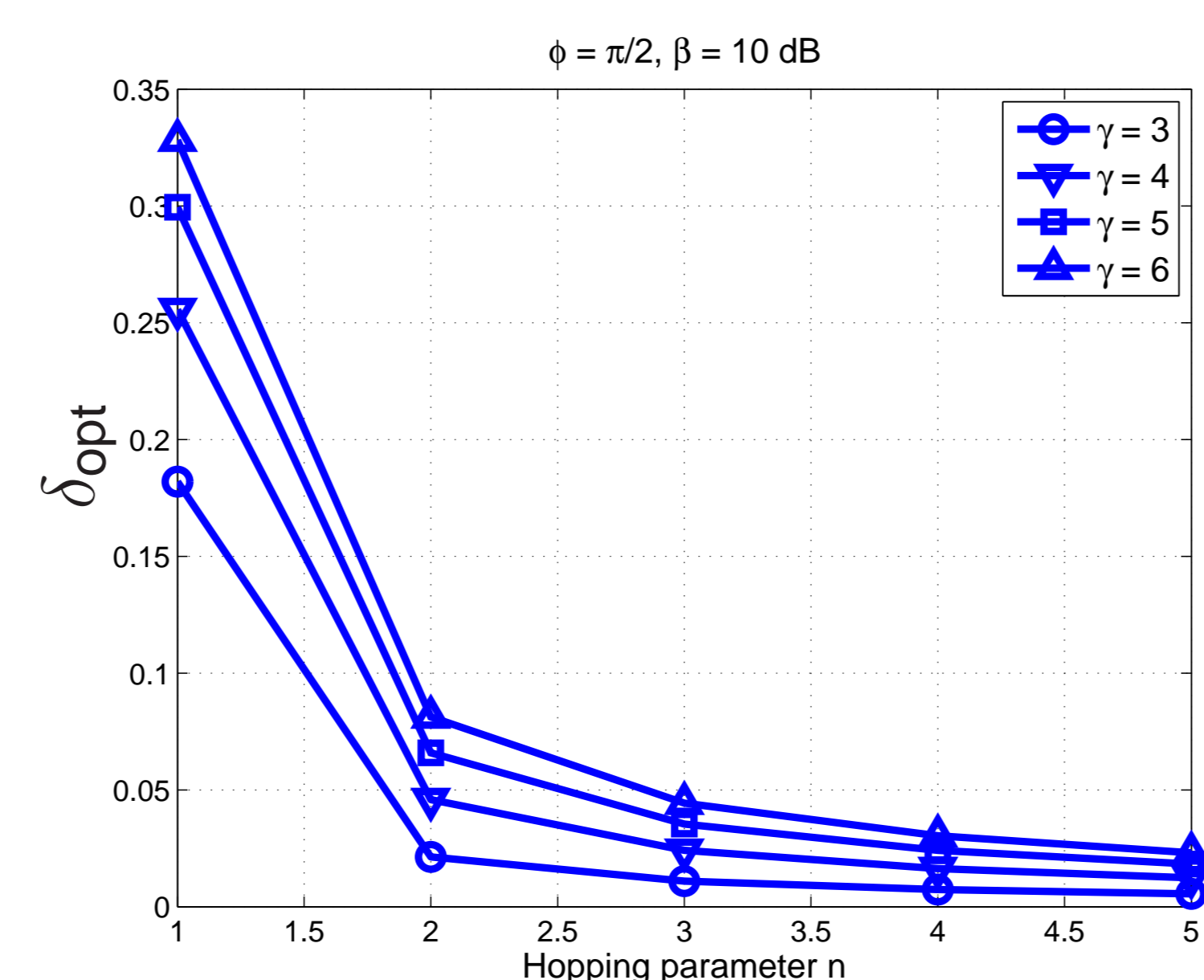
where $c = \pi\Gamma(1 + 2/\gamma)\Gamma(1 - 2/\gamma)\beta^{2/\gamma} = \frac{2\pi^2\beta^{2/\gamma}}{\gamma \sin(2\pi/\gamma)}$.

Proposition 2 For the r-TDMA-based flow with N relays and success probability p_s , the throughput at steady state is

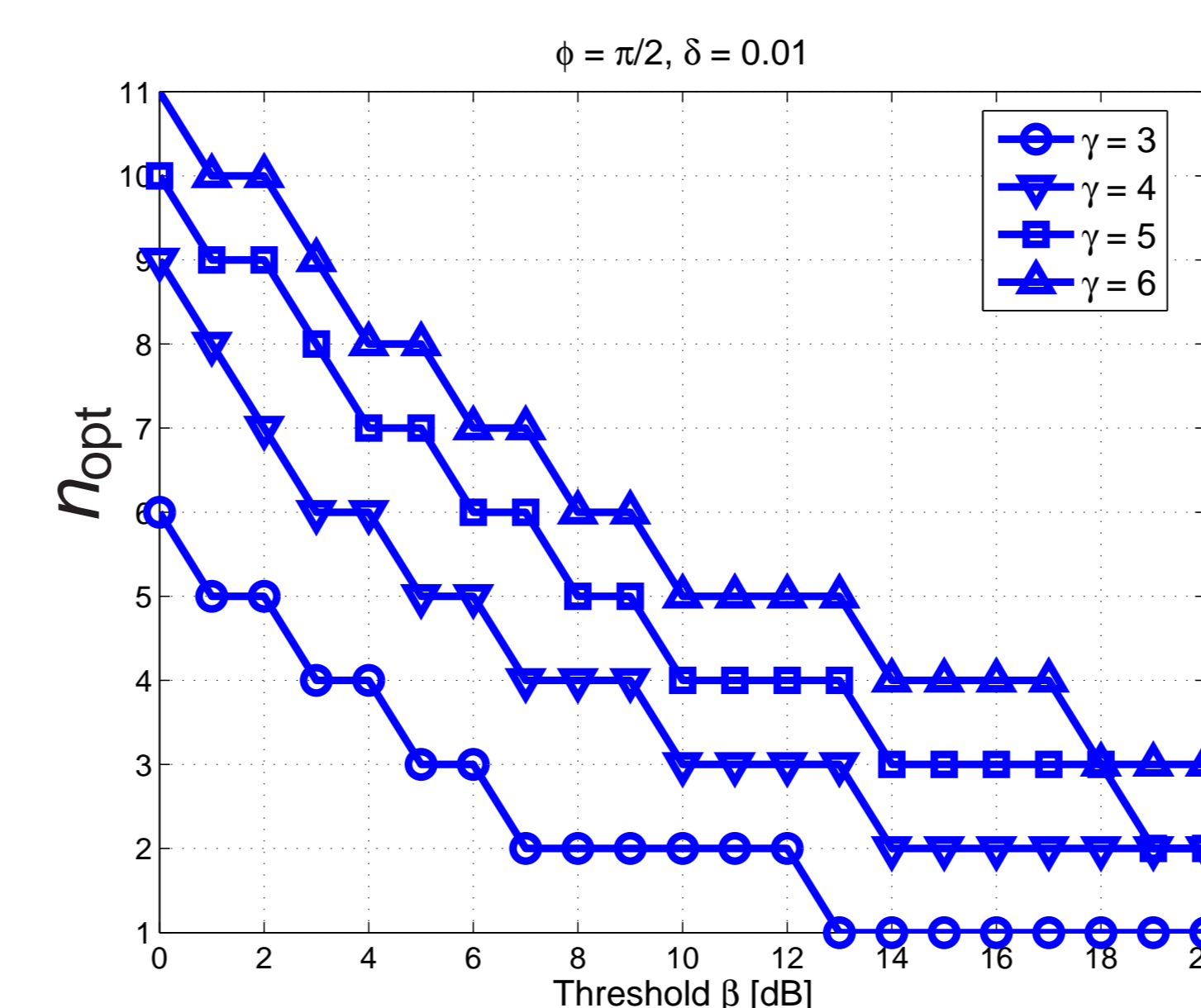
$$T = \frac{p_s}{2N + 1}.$$

Using the definition of transport density, we obtain

$$\rho_{\text{trans}} = \frac{N}{(2N + \sqrt{n})} \frac{\delta(1-\delta)^{n-1} \phi^{n-3/2}}{((1-\delta)\phi + 2\delta c)^n} \sqrt{2\pi n} \sin\left(\frac{\phi}{2}\right).$$



$$\delta_{\text{opt}} = \frac{3\phi + 4(n-1)c - \sqrt{(\phi + 4(n-1)c)^2 + 16n\phi c}}{2(\phi - 2c)}.$$



$$n_{\text{opt}} \approx \left\lceil 0.5 / \ln\left(1 + \frac{2\delta c}{(1-\delta)\phi}\right) \right\rceil.$$

ANALYSIS OF ALOHA-BASED WIRELESS NETWORKS

For analytical tractability, we assume the following.

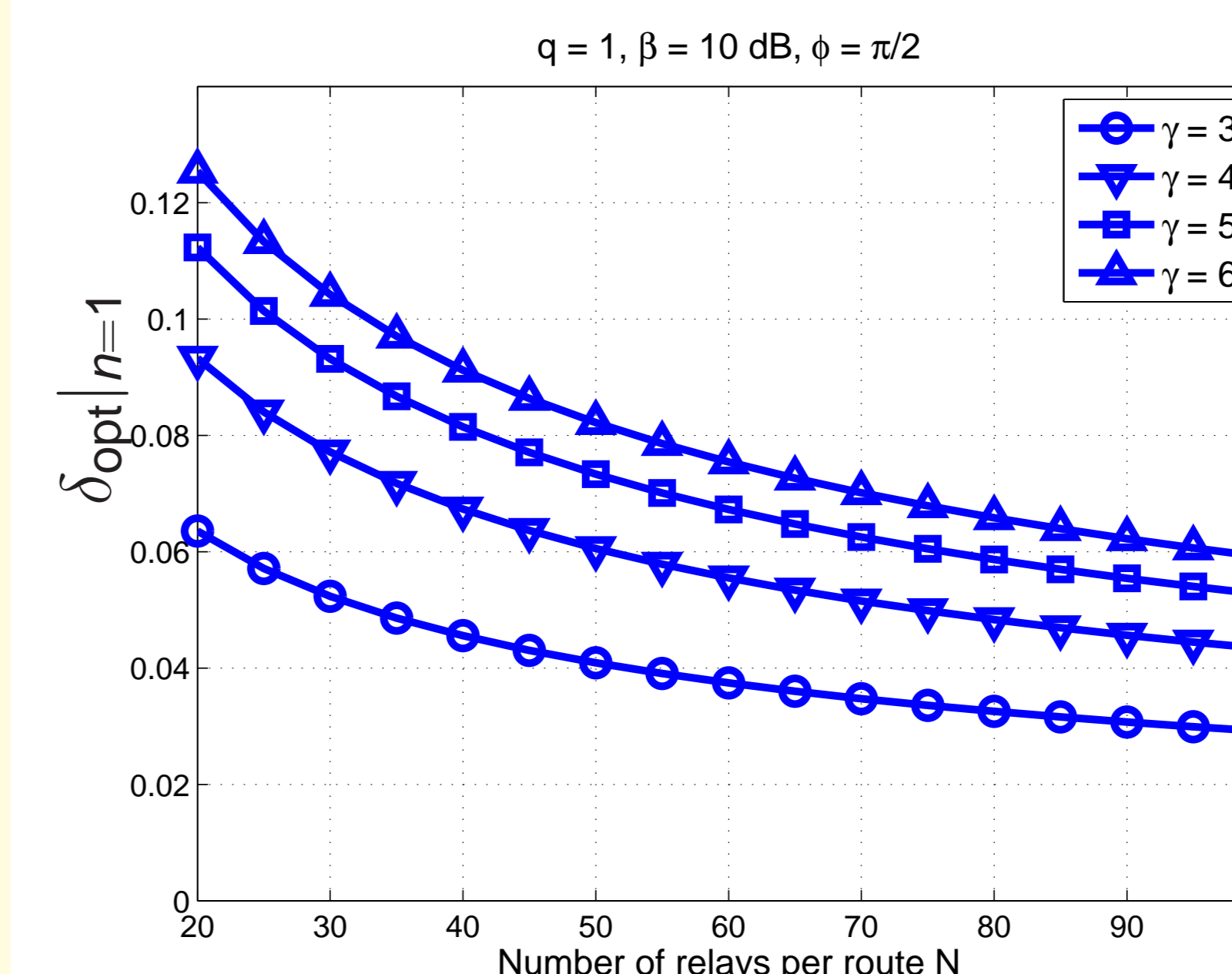
- We consider long flows, i.e., take $N \gg 1$.
- We employ the *mean field theory* assumption.

Proposition 3 For the ALOHA-based network with N relay nodes, the packet success probability p_s from any node to its n^{th} -nearest-neighbor is

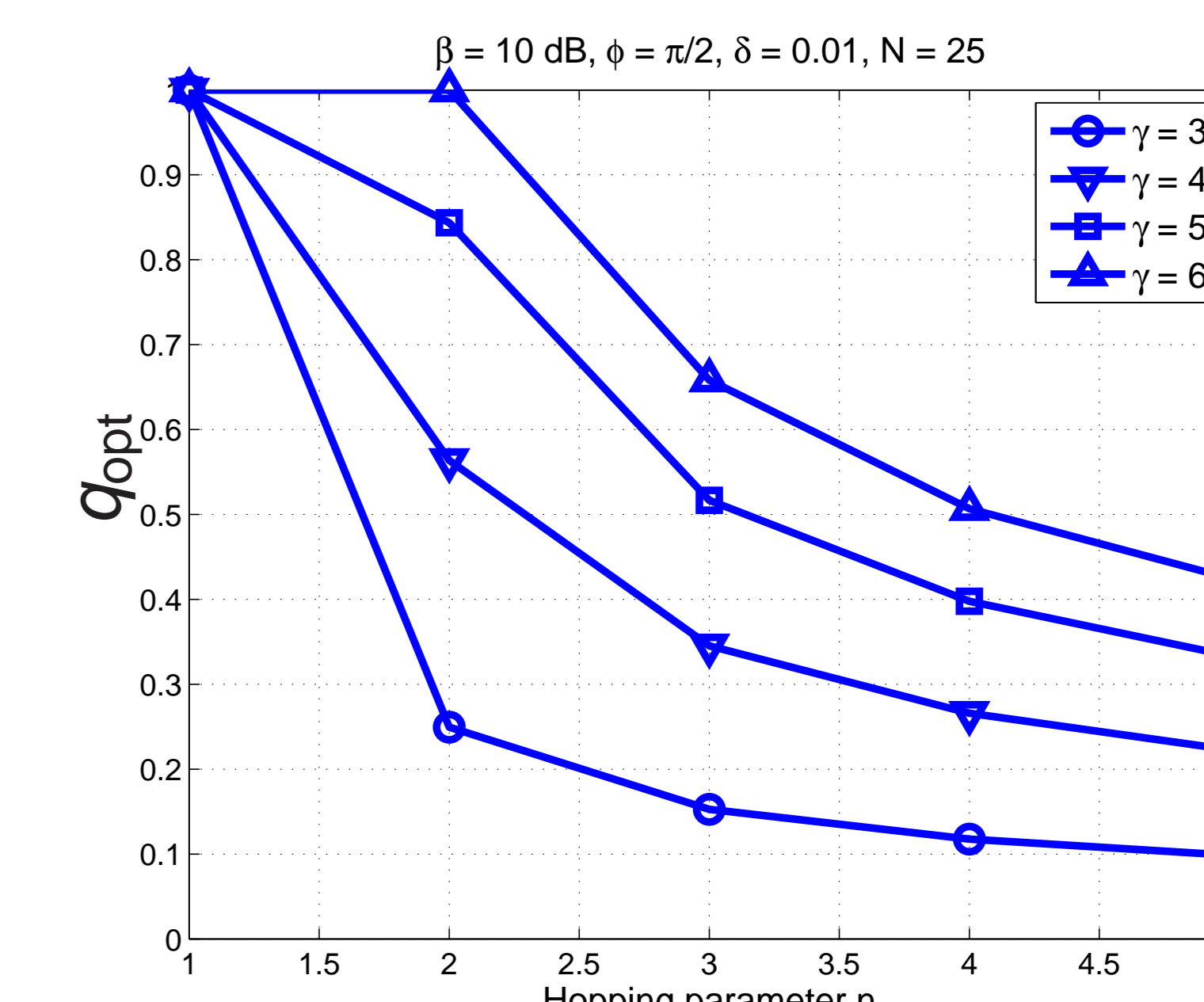
$$p_s \approx \left(\frac{(1-\delta)\phi}{(1-\delta)\phi + \delta q N c} \right)^n.$$

The transport density may be lower-bounded as

$$\rho_{\text{trans}} \approx \frac{Nq\delta(1-\delta)^{n-1} \phi^{n-3/2} n^{n/2}}{((1-\delta)\phi + \delta q N c)^n} \sqrt{\frac{\pi}{8}} \sin\left(\frac{\phi}{2}\right).$$

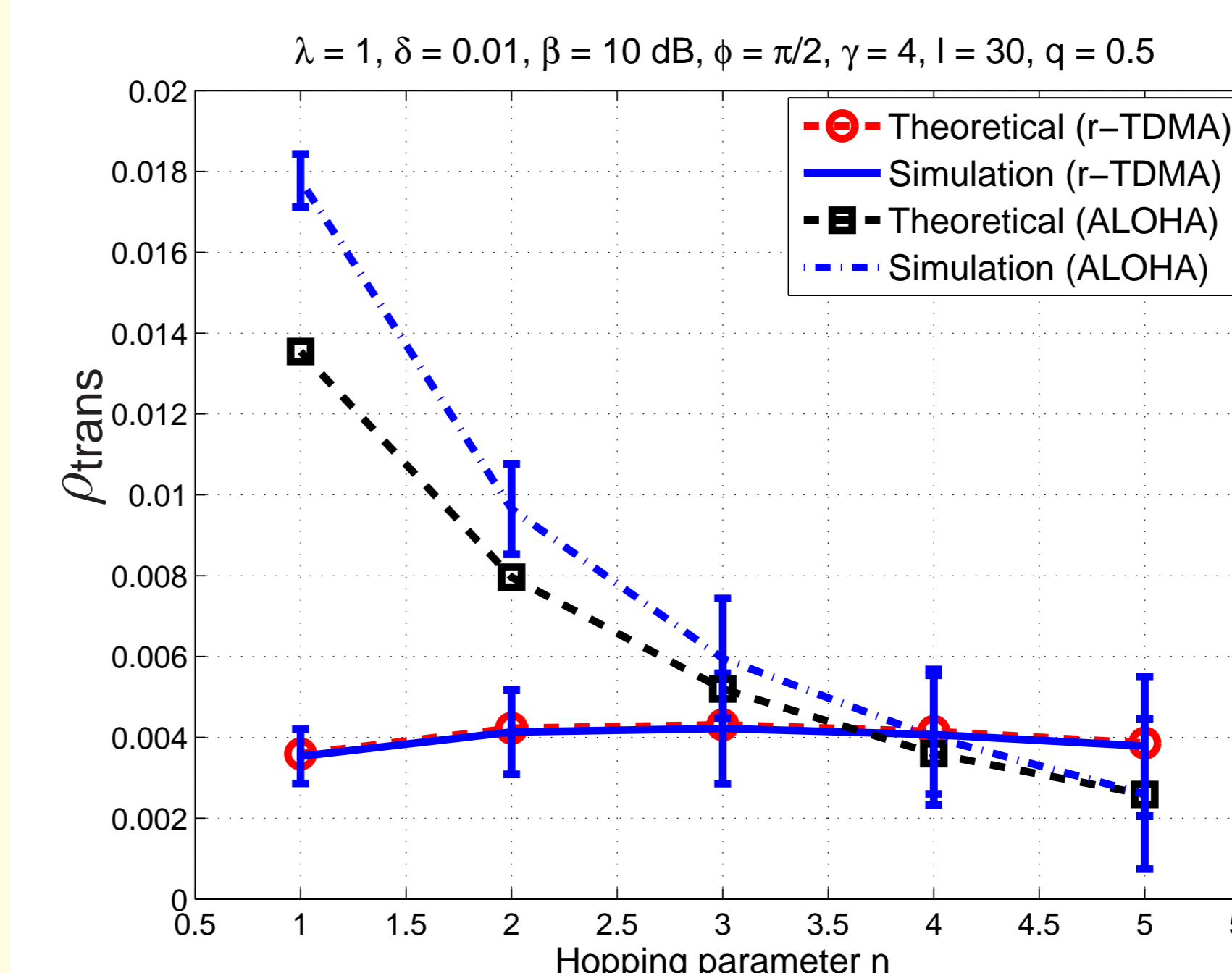


$$\delta_{\text{opt}}|_{n=1} = \frac{3\phi - \sqrt{\phi^2 + 8\phi N c}}{2(\phi - N c)}.$$



$$q_{\text{opt}} = \min\left\{1, \frac{(1-\delta)\phi\sqrt{n}}{\delta N c(n-1)}\right\}.$$

SIMULATION RESULTS



- Verification needed since relays may serve multiple nodes.
- Based on 100 different realizations of the point process on a 40×40 square.
- There exist regimes where either scheme performs better.

FUTURE WORK

- Propose distributed routing algorithms for MANETs that are 'optimal' in nature.
- Characterize the performance of other MAC schemes, in particular CSMA and spatial TDMA.
- Use ideas from *Langmuir Kinetics* to characterize the TDR region in wireless ad hoc networks.